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## RESEARCH ARTICLE

# Attcack and Saving Network from Attack

## \*Luaay Abdulwahed

College of nursing, Basra University, Iraq \*Corresponding author: Loea1965@yahoo.com, luaay.shihab@uobasrah.edu.iq

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## **ABSTRACT**

This paper takes a keen look in network security and network security goals, why computers are insecure and type Of attack, how to improve network security. Securing your network and practical's what is Quota.

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## INTRODUCTION

Network security comprises the measures a company takes to protect its computer system, and it is a prime concern for every company that uses computers. Compromised network security means a hacker or competitor may gain access to critical or sensitive data, possibly resulting in data loss, or even complete destruction of the system.

## **Network Security Goals**

- Confidentiality: only sender, intended receiver should "understand" message contents sender encrypts message receiver decrypts message Privacy: hide 'who is doing what with whom
- Authentication: sender, receiver want to conform identity of each other
- Integrity: sender, receiver want to ensure messages are not altered (in transit, or afterwards) without detection

Access and Availability: services must be accessible and available to users Network Security

## Why Computers are Insecure?

## Most PCs use insecure Ross

- Most designed for 'home' security not a goal
- Others support separation between users
- Few/none restrict capabilities of applications
- Malicious / vulnerable / buggy app can harm all!!

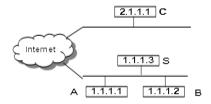
PCs run buggy, vulnerable, even malicious code Many sources (libraries, shareware, ...)

Limited awareness & tools.

Limited product liability and consequent damages Most computers don't fix known vulnerabilities

# Security Flaws in IP

- The IP addresses are filled in by the originating host
  - Address spoofing
- · Using source address for authentication
  - r-utilities (rlogin, rsh, hosts etc..)



·Can A claim it is B to the server ? •ARP Spoofing •Can C claim it is B to the server 8? ·Source Routing

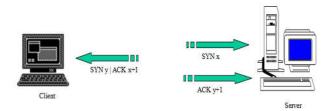
## Routing Attacks

- Distance Vector Routing

   Announce 0 distance to all other nodes
  - · Black hole traffic
  - Eavesdrop
- Link State Routing
  - Can drop links randomly
  - Can claim direct link to any other routers A bit harder to attack than DV

  - BGP
    - ASes can announce arbitrary prefix
    - ASes can alter path

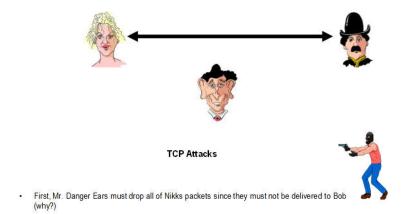
### TCP Attacks

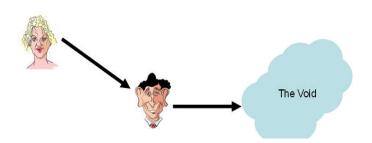


- Server needs to keep waiting for ACK y+1 Server recognizes Client based on IP address/port and y+1

## TCP Attacks

Nikks and Luaay have an established TCP connection





### **TCP Attacks**

Why are these types of TCP attacks so dangerous?







Trusting web client

## TCP Layer Attacks

### TCP Session Hijack

- When is a TCP packet valid?
  - · Address/Port/Sequence Number in window
- How to get sequence number?
  - · Sniff traffic
  - Guess it
    - Many earlier systems had predictable ISN
- Inject arbitrary data to the connection

## TCP Session Poisoning

- Send RST packet
  - · Willtear down connection
- Do you have to guess the exact sequence number?
  - Anywhere in window is fine
  - · For 64k window it takes 64k packets to reset
  - · About 15 seconds for a T1

## Attacks on UDP

- Fields in first 64 bits: Source Port, Destination Port, Length, Checksum Bits 65-128: Data payload
- Dest. Port: Decrypted packets delivered to X<sub>B</sub>.
- Length: Packets can be truncated.
- Checksum can be fixed directly
- With a 128-bit cipher, the first 64 bits of the payload can be modified.

## IV Attacks on L2TP

Layer 2 Tunneling Protocol (L2TP):

- Used to tunnel PPP connections over a wide area network. L2TP data packets encapsulated in UDP
- Typically uses IPsec ESP for encryption

IV attack on L2TP with IPsec ESP:

- 128-bit IV reaches the first 64 bits of L2TP packet
- contains Length, Tunnel ID, Call ID
- Call ID attack: Decrypted packet can be delivered to attacker's connection. (similar to attack on TCP/UDP Destination Port)

## Attacks On Application Layer

Applications don't authenticate properly Authentication information in clear FTP, Telnet, POP

## DNS insecurity

- DNS poisoning
- DNS zone transfer

### Types of attacks

- Denial of Service (DoS) Attacks
- □□ Website Defacement
- □□ Viruses and Worms
- □□ Data sniffing and Spoofing
- □□ Unauthorized Access
- □□ Malicious Code and Trojans
- □□ Port-scanning and Probing
- □□ Wireless Attacks

### Denial Of Service (Clogging) Attack

- Denial Of Service (Clogging) Attack Attacker tries to exhaust resources of host /server / router / user
- Resources include:-
- Computations (CPU time) Storage
- (e.g. for state of requests/connections)
  Open TCP connections
- Limited (10s to several thousand connections -depending on hardware, operating system)
- \_\_\_\_\_observer in unousand connections -depending on hardware, operating system So server inever keeps open connections !Always request-response (and server closes connection, no state)
- SYN flooding DOS attack: attacker sends' SY N' flow (open connection); server waits... Network Security

## SYN flooding DOS (clogging) Attack

- Recall T CP connection setup process
- Attacker sends many SY N requests (using different spoofed client IP addresses), no ACK
- Uses up server's capacity for open connections
- Possible solution: request must contain 'cookie' (next)

## Attacks Through The Net

- Eavesdropping
- Port scanning (probing for weaknesses)
- Spoofing-
  - Fake e-mail
  - Using a fake IP address
- Denial of Service (DoS)-

Shut down the target host via a critical fault-Also available in distributed format to simply overload a target Message replay

Connection capture (TCP)

## Port-scanning

- Technique that identifies vulnerable network ports or services (i.e. TELNET, FTP, E-mail, Web, etc)
- Works by identifying as many targets as possible and tracking the ones that are receptive
- Scanning software is free and commonly accessible via the web

## (Spoofing)

- IP spoofing An attacker may fake their IP address so the receiver thinks it is sent from a location that it is not actually from. There are various forms and results to this attack.
- The attack may be directed to a specific computer addressed as though it is from that same computer. This may make the computer think that it is talking to itself. This may cause some

operating systems such as Windows to crash or lock up.

#### Man in the middle attack -

Session hijacking - An attacker may watch a session open on a network. Once authentication is complete, they may attack the client computer to disable it, and use IP spoofing to claim to be the client who was just authenticated and steal the session. This attack can be prevented if the two legitimate systems share a secret which is checked periodically during the session.

#### (DNS POISONING)

- This is an attack where DNS information is falsified. This attack can succeed under the right conditions, but may not be real practical as an attack form. The attacker will send incorrect DNS information which can cause traffic to be diverted. The DNS information can be falsified since name servers do not verify the source of a DNS reply. When a DNS request is sent, an attacker can send a false DNS reply with additional bogus information which the requesting DNS server may cache. This attack can be used to divert users from a correct web server such as a bank and capture information from customers when they attempt to logon.

#### Some DoS Attacks

- **Ping broadcast** . A ping request packet is sent to a broadcast network address where there are many hosts. The source address is shown in the packet to be the IP address of the computer to be attacked. If the router to the network passes the ping broadcast, all computers on the network will respond with a ping reply to the attacked system. The attacked system will be flooded with ping responses which will cause it to be unable to operate on the network for some time, and may even cause it to lock up. The attacked computer may be on someone else's network. One countermeasure to this attack is to block incoming traffic that is sent to a broadcast
- Ping of death An oversized ICMP datagram can crash IP devices that were made before 1996
- Smurf An attack where a ping request is sent to a broadcast network address with the sending address spoofed so many ping replies will come back to the victim and overload the ability of the victim to process the replies.
- Teardrop a normal packet is sent. A second packet is sent which has a fragmentation offset claiming to be inside the first fragment. This second fragment is too small to even extend outside the first fragment. This may cause an unexpected error condition to occur on the victim host which can cause a buffer overflow and possible system crash on many operating systems.

## Wireless Attacks

- Wireless Equivalent Privacy (WEP) protocol cannot be trusted for security
- Attackers can easily eavesdrop or spoof wireless traffic
- Hackers external to your building may be able to intercept and view all of your wireless traffic, despite encryption
- Hacker tools free and easily accessible via the web: AirSnort, WEPCrack.

## How to improve network security

- General awareness of Network Security among users Upgrading the skill of the system and network administrators
- Sharing of network security information, knowledge and experience amongst system and network administrators

### Countermeasure

- Personnel Security Policy and Procedures
- Training and Awareness
- Physical Security
- Dedicated Management Technology
- Firewalls
- Intrusion Detection Virus Protection
- Authentication and Authorization
- Encryption
- Auditing and Assessment (Third Party)
  Data and Information Backup

### SSH

SSH can be used for secured Command and Control sessions to routers.

- · Full SSH has three components
- · a terminal session with a secure transport
- the ability to handle "r-commands" similar to rsh
- the ability to "forward" other TCP-based protocols

### **SSH Authentication**

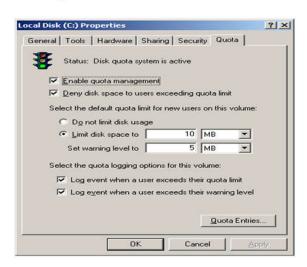
- · There are two levels of
- · Authentication required for an SSH session
- · Host (or 'device') Authentication
- User Authentication

### What is Quota

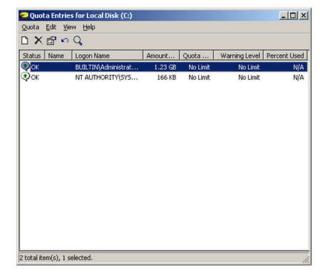
 To select other options as desired: enable disk quotas for the users, log on as an administrator, right click on the drive you want to limit, and select the quota tab. Check Enable quota management, and select other options as desired:

## Quota Management

- Open properties of any drive on which you want to add quota limit.
- Shown in figure 1.1



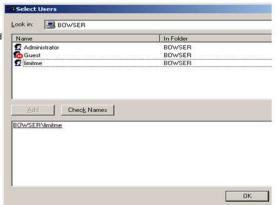
- Then select Quota tab and select new entry to add quota limit.
- Shown in figure 1.2



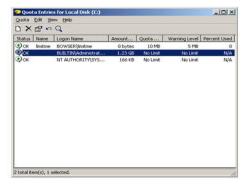
- Then select space which you want to provide a user for the use..
- Shown in figure 1.3



- Then select user on which you want to add quota limit.
- Shown In figure 1.4



- After selecting user you will find that your user is in the list on which you r applying quota limit. And then click finish.
- Shown in figure 1.5



- After applying quota limit a user will find this error log.
- Shown in figure 1.6



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